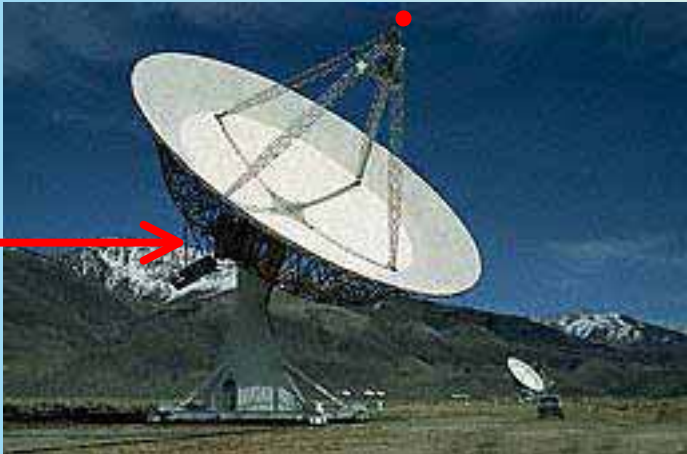


Motivation: Classical Program

Classical program : sequence of bits



01101

Motivation: Quantum Program

Quantum program
State vector of n qubits



2 bits of classical information



EPR



Motivation: Quantum Program

Quantum program
State vector of n qubits



2 bits of classical information



EPR



Universality ?



Research Center for Quantum Information
Bratislava, Slovakia

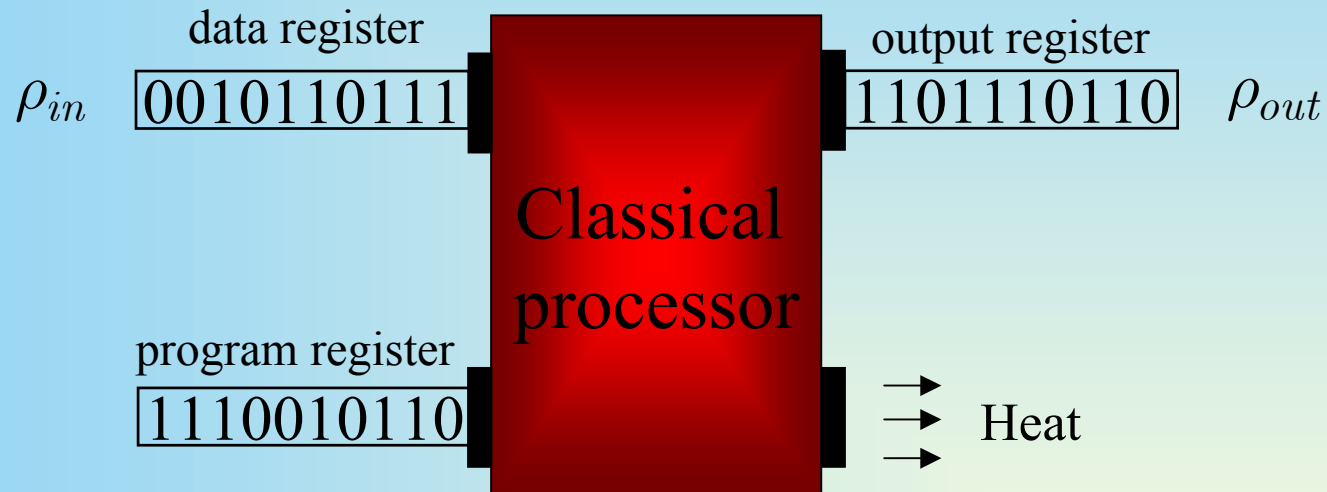
Programmable Quantum Processors

Vladimír Bužek
Mark Hillery
Mário Ziman

QUBITS → QGATES

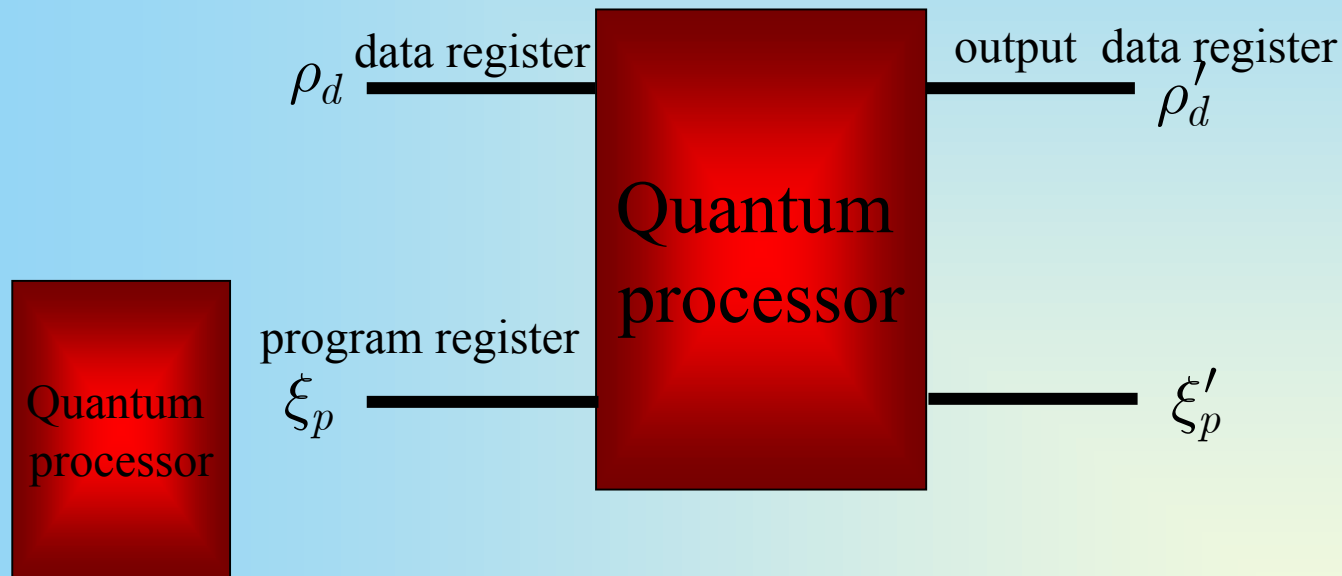
16 September, 2002, Dublin

Model of Classical Processor



$$\rho_{out} = T[\rho_{in}]$$

Quantum Processor



Quantum processor – **fixed** unitary transformation U_{dp}

\mathcal{H}_d – data system,

$S(\mathcal{H}_d)$ – data states

\mathcal{H}_p – program system,

$S(\mathcal{H}_p)$ – program states

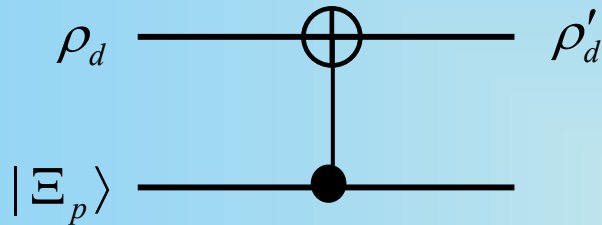
Two Scenarios

- **Measurement-based strategy** - estimate the state of program

$$F = \frac{N + 1}{N + d}$$

- **Quantum strategy** – use the quantum program register
conditional (probabilistic) processors
unconditional processors

C-NOT as Unconditional Quantum Processor



$$\text{CNOT } |\psi\rangle|0\rangle = |\psi\rangle|0\rangle$$

$$\text{CNOT } |\psi\rangle|1\rangle = \sigma_x |\psi\rangle \otimes |1\rangle$$

- program state $|0\rangle \Rightarrow \mathbf{1}$ implemented, i.e. $\rho_d \rightarrow \rho'_d = \rho_d$
- program state $|1\rangle \Rightarrow \sigma_x$ implemented, i.e. $\rho_d \rightarrow \rho'_d = \sigma_x \rho_d \sigma_x$
- general pure state $|\Xi_p\rangle = \alpha|0\rangle_p + \beta|1\rangle_p \Rightarrow \rho_d \mapsto \rho'_d = |\alpha|^2 \rho_d + |\beta|^2 \sigma_x \rho_d \sigma_x$
- unital operation, since $\Phi[\mathbf{1}] = |\alpha|^2 \mathbf{1} + |\beta|^2 \sigma_x \mathbf{1} \sigma_x = \mathbf{1}$
- program state is **2-d** and we can apply **2** unitary operations

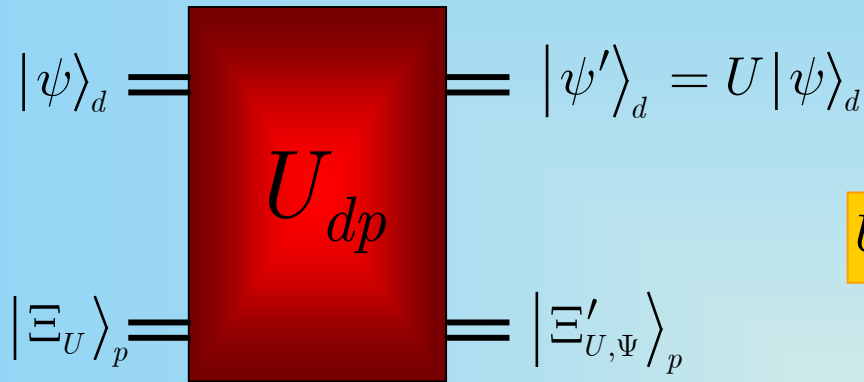


Question (Nielsen & Chuang)

Is it possible to build a *universal* programmable quantum gate array which take as input a quantum state specifying a quantum program and a data register to which the unitary operation is applied ?

on a qubit an ∞ number of operations can be performed

Theorem



$$U_{dp} (|\psi\rangle_d \otimes |\Xi_U\rangle_p) = (U|\psi\rangle) \otimes |\Xi'_{U,\psi}\rangle$$

- **no** universal deterministic quantum array of **finite** extent can be realized
- on the other hand – a program register with ***d*** dimensions can be used to implement ***d*** unitary operations by performing an appropriate sequence of controlled unitary operations

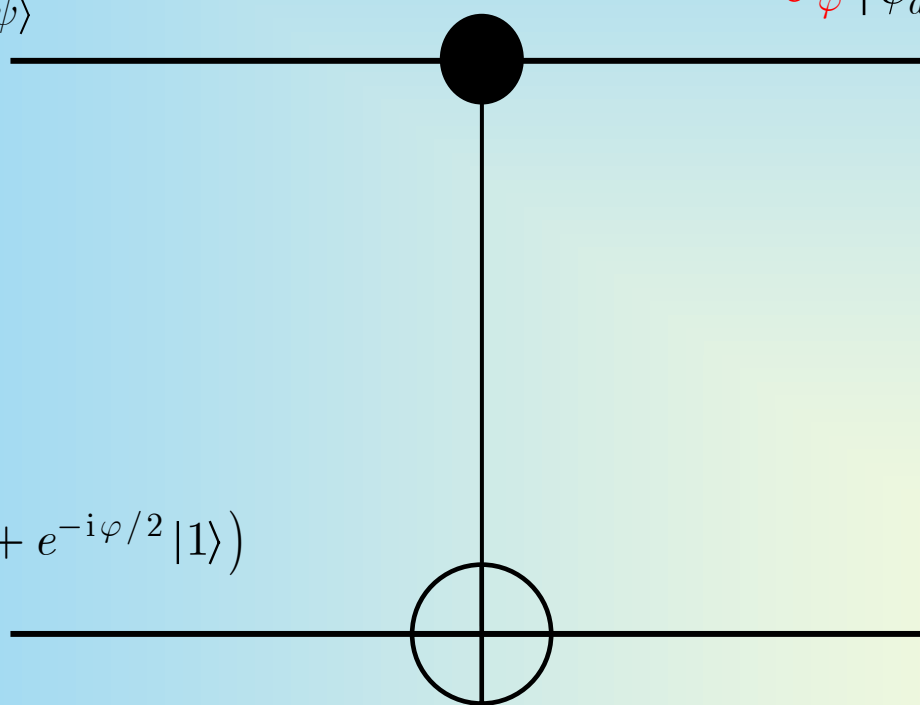
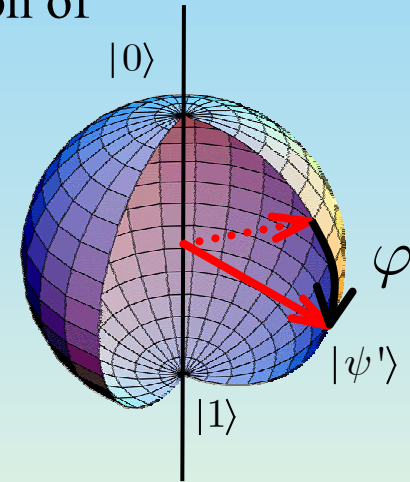
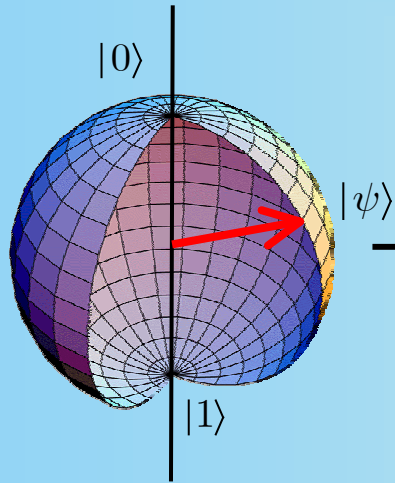
$$|\text{answer}\rangle = |\text{NO}\rangle|\text{deterministic}\rangle + |\text{YES}\rangle|\text{probabilistic}\rangle$$

C-NOT as Probabilistic Quantum Processor

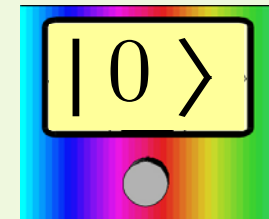
- Vidal & Cirac – probabilistic implementation of

$$U_\varphi = e^{-i\frac{\varphi}{2}\sigma_z}$$

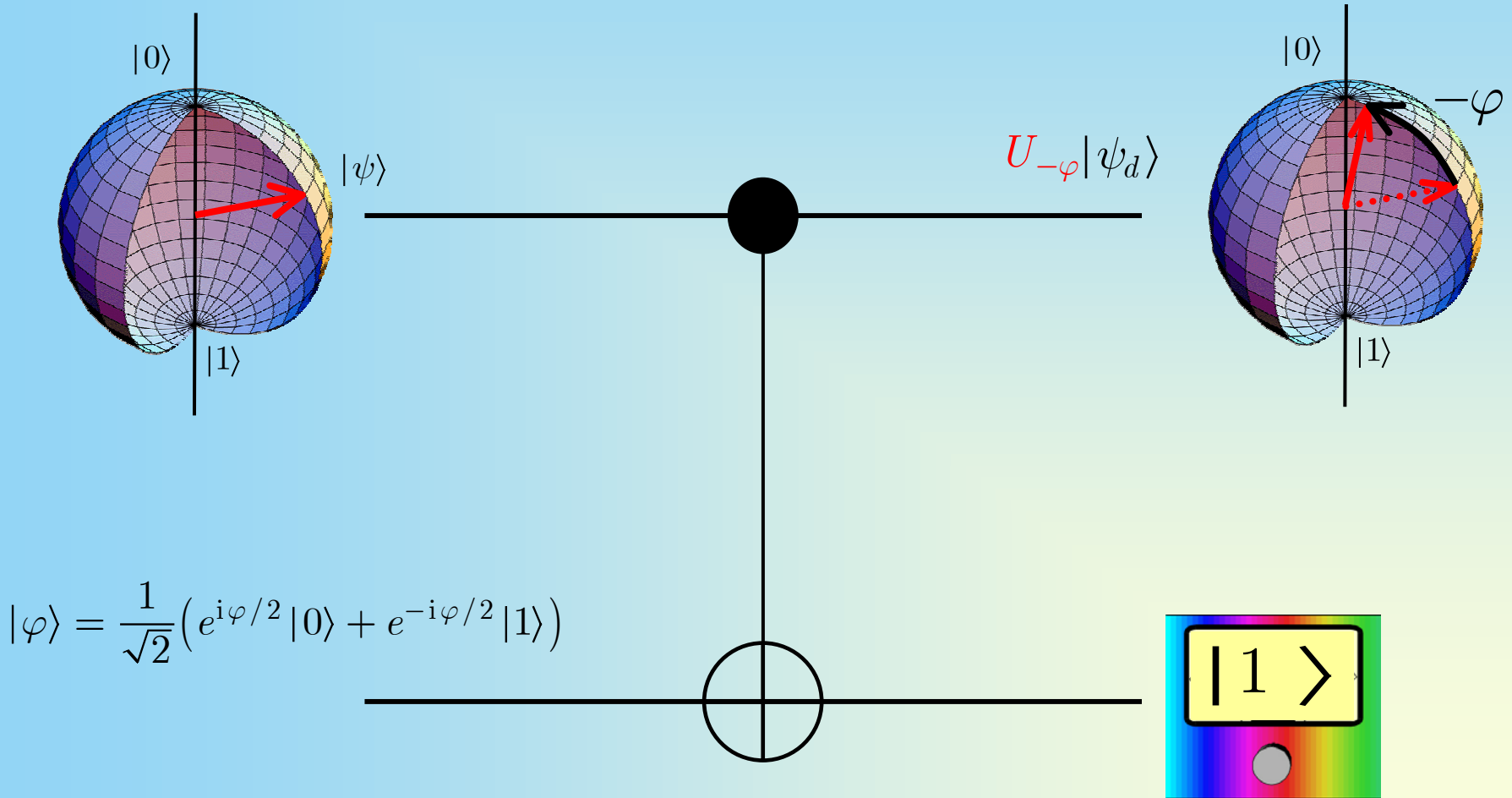
$$U_\varphi |\psi_d\rangle$$



$$|\varphi\rangle = \frac{1}{\sqrt{2}}(e^{i\varphi/2}|0\rangle + e^{-i\varphi/2}|1\rangle)$$

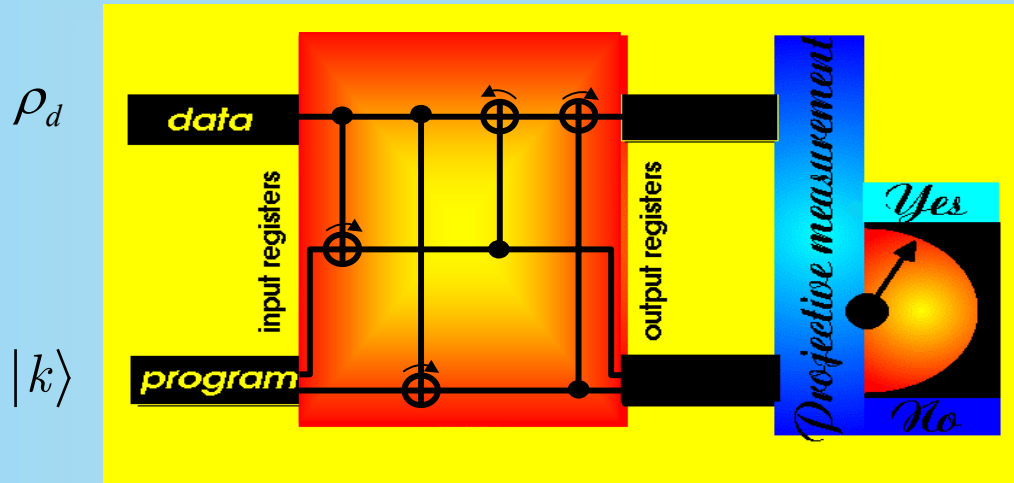


C-NOT as Probabilistic Quantum Processor



Correction of the error – new run of the processor with $|2\varphi\rangle$

Universal Probabilistic Processor



ρ_d

$$|\xi_p\rangle = \sum_k \alpha_k |k\rangle$$

$$\rho'_d = A\rho_d A^\dagger$$

$$A = \sum_k \alpha_k U_k$$

$$|M\rangle = \frac{1}{D} \sum_k |k\rangle$$

- implementation of any unitary transformation
- **BUT!** only with probability and $P_{\text{success}} = (\dim \mathcal{H}_p)^{-1}$
- Example (Quantum information distributor)

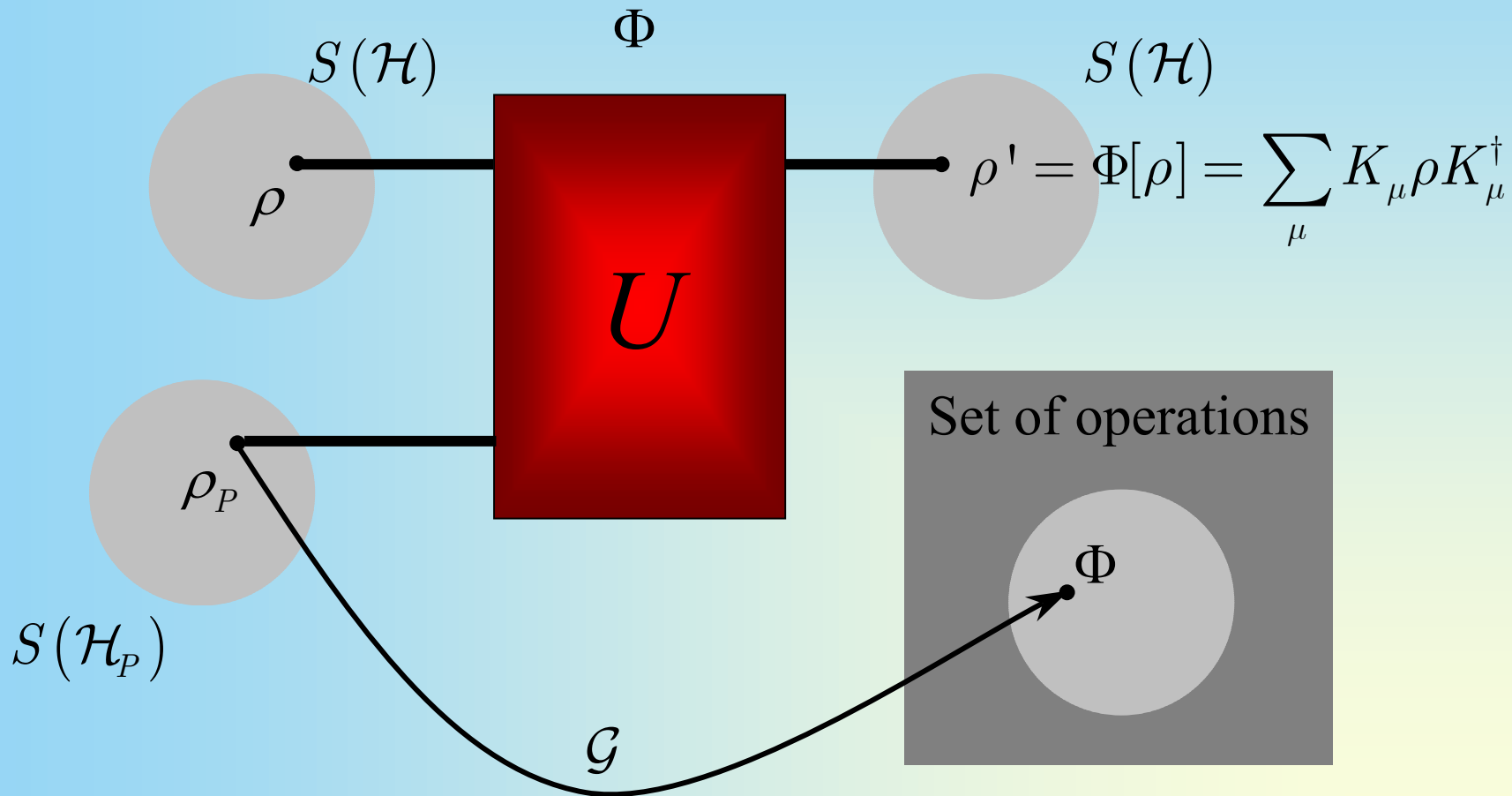
$$\mathcal{H}_p = \mathcal{H}_d \otimes \mathcal{H}_d ; D = \dim \mathcal{H}_d$$

$\{|k\rangle\}_k = \{|k'\rangle\}_k \dots$ basis of maximally entangled states of \mathcal{H}_p

$\{U_k\}_k \dots$ set of D^2 unitaries forming an operator basis, $\text{Tr } U_k^\dagger U_l = D\delta_{kl}$

- subset of unitary transformations – higher probability

Implementation of Maps via Unconditional Quantum Processors



Mathematical Description of Quantum Processors

- definition of U_{dp} via “Kraus operators” $A_{kl} := {}_p \langle l | U_{dp} | k \rangle_p$

$$U_{dp} \left(|\psi\rangle_d \otimes |k\rangle_p \right) = \sum_l \left(A_{kl} |\psi\rangle_d \right) \otimes |l\rangle_p$$

- normalization condition $\sum_l A_{k_1 l}^\dagger A_{k_2 l} = \delta_{k_1 k_2} \mathbf{1}_d$
- induced quantum operation $\rho_d \mapsto \rho'_d = \Phi_k [\rho_d] = \sum_l A_{kl} \rho_d A_{kl}^\dagger$
- general pure program state $|\Xi\rangle_p = \sum_k \alpha_k |k\rangle_p$

$$\rho_d \mapsto \rho'_d = \Phi_\Xi [\rho_d] = \sum_l A_l(\Xi) \rho_d A_l^\dagger(\Xi)$$

$$A_l(\Xi) = {}_p \langle l | U_{dp} | \Xi \rangle_p = \sum_k \alpha_k A_{kl}$$

- can be generalized for mixed program states

Classes of Processors

- Induced mapping

\mathcal{G} : program states \rightarrow set of operations

$$\xi_p \rightarrow \mathcal{G}(\xi_p) \equiv \Phi_\xi \quad ; \quad \Phi_\xi : S(\mathcal{H}_d) \rightarrow S(\mathcal{H}_d)$$

- Classes of processors

1. U processors – implement unitary operations
2. covariant processors
3. injective processors – each ξ_p encodes different operation Φ_ξ

U processors

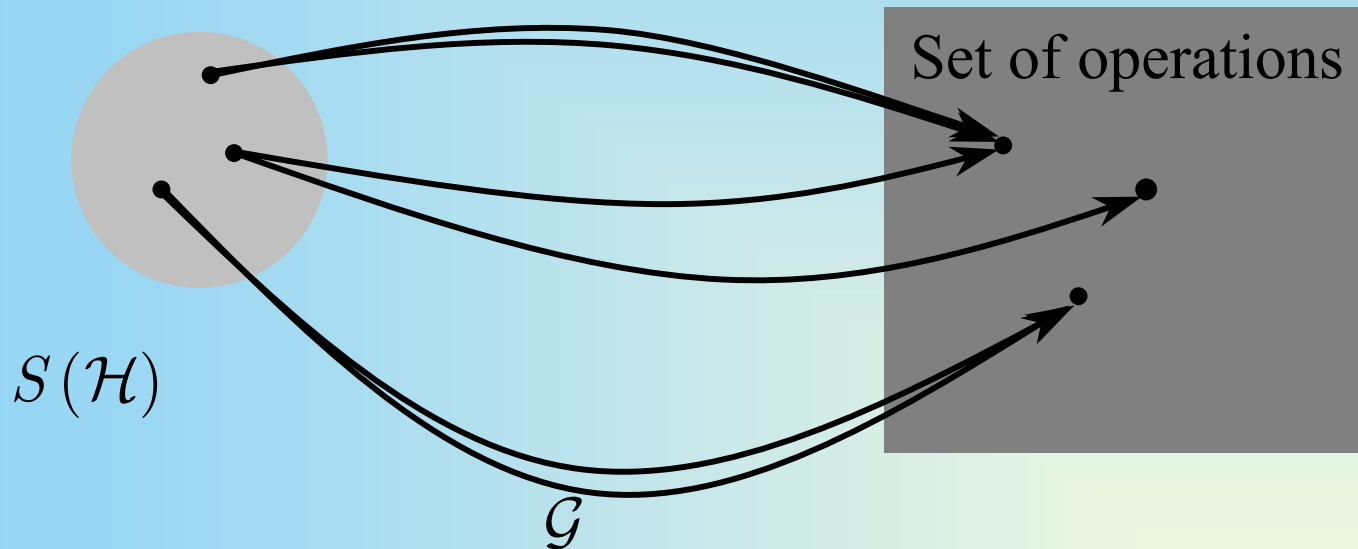
$$U_{dp} \left(|\psi\rangle_d \otimes |k\rangle_p \right) := \left(U_k |\psi\rangle_d \otimes |k'\rangle_p \right)$$

the most general processor that implements $N = \dim \mathcal{H}_p$
different unitaries

elementary programs $|k\rangle_p$ realize unitary transformation U_k
general program state ξ_p implements **unital operations**, i.e. $\Phi[\mathbf{1}] = \mathbf{1}$

Injective Processors

the induced mapping \mathcal{G} is injective, i.e. \forall pair $\Xi_{1p} \neq \Xi_{2p} \Rightarrow \Phi_1 \neq \Phi_2$



Does such processor exist? **YES** partial swap processors:

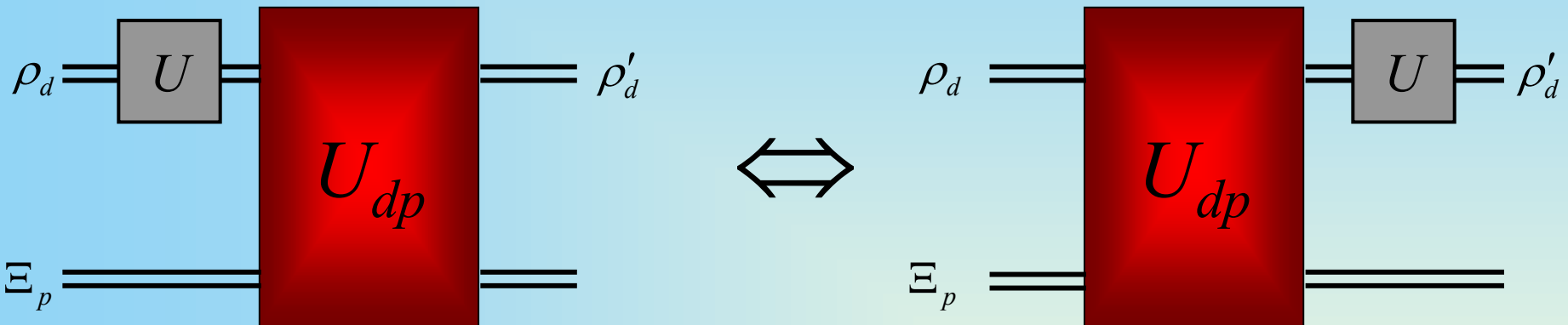
$$P_\theta = \cos \theta \mathbf{1} + \sin \theta S \quad S(|\psi\rangle \otimes |\phi\rangle) = |\phi\rangle \otimes |\psi\rangle, \quad \dim \mathcal{H}_p = \dim \mathcal{H}_d$$

program Ξ_p encodes contractive operation Φ_Ξ with the only fixed point Ξ_d

$$\Phi_\Xi[\Xi_d] = \Xi_d$$

Covariant Processors

- the covariance condition



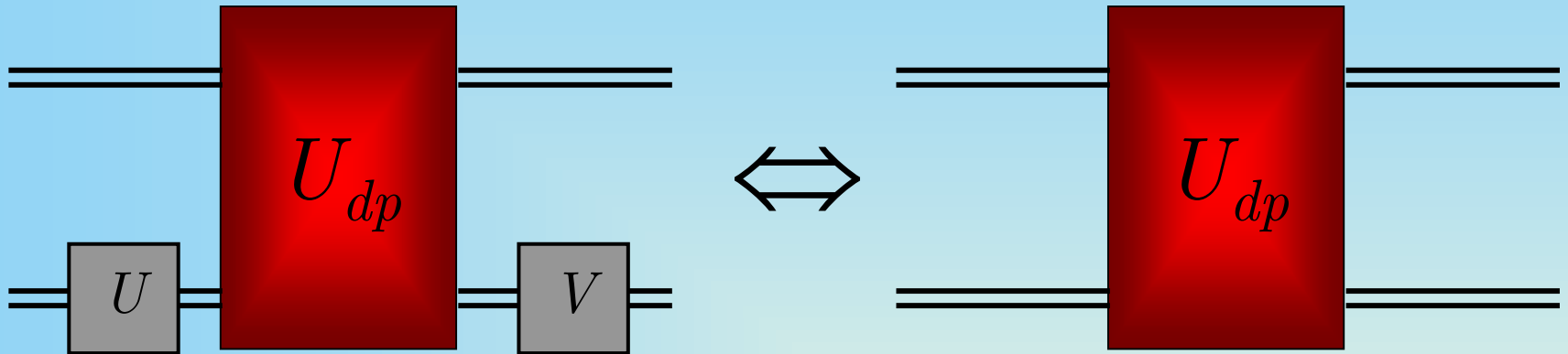
$$\Phi_{\Xi} [U \rho_d U^{\dagger}] = U \Phi_{\Xi} [\rho_d] U^{\dagger}$$

$$\forall \Xi \in S \subset S(H_p)$$

$$\forall U \in SU(\dim H_d)$$

- covariance with respect to the program states from the set S
- if $S = S(H)$, then the existence of nontrivial covariant U_{dp} is questionable
- if U_{dp} implements unitary transformation \Rightarrow it is not covariant in this sense

Equivalent Processors



if $U_{dp}^{(2)} = (\mathbf{1} \otimes V) U_{dp}^{(1)} (\mathbf{1} \otimes U)$, then $U_{dp}^{(1)} \sim U_{dp}^{(2)}$

- equivalent processors implement the same quantum operations
- basis operators $A_{jk}^{(1)}$ and $A_{jk}^{(2)}$ associated with the two processors satisfy

$$A_{jk}^{(2)} = \sum_{m,n=1}^N (U_{p1})_{jm} (U_{p2})_{nk} A_{mn}^{(1)}$$



Inverse Problem

“Given a set Φ_x of quantum operations . Is it possible to design a processor that performs all these operations?”

1. Continuous set of unitaries = question of universality **NO**
2. Phase damping channel = model of decoherence **YES**
3. Amplitude damping channel = model exponential decay **NO**



Conclusions & Open Questions

- programmable quantum computer – programs via quantum states programs can be outputs of another QC
- some CP maps via unconditional quantum processors
- arbitrary CP maps via probabilistic programming
- controlled information distribution (eavesdropping)
- simulation of quantum dynamics of open systems
- set of maps induced by a given processor (loops)
- quantum processor for a given set of maps
- quantum multi-meters

M.Hillery, V.Buzek, and M.Ziman: *Phys. Rev. A* 65, 022301 (2002).

M.Dusek and V.Buzek: *Phys. Rev. A* 66, 022112 (2002).

M.Hillery, M.Ziman, and V.Buzek: *Phys. Rev. A* 66, 042302 (2002)